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# CISC 235: Topic 11

## Shortest Paths Algorithms

# Outline

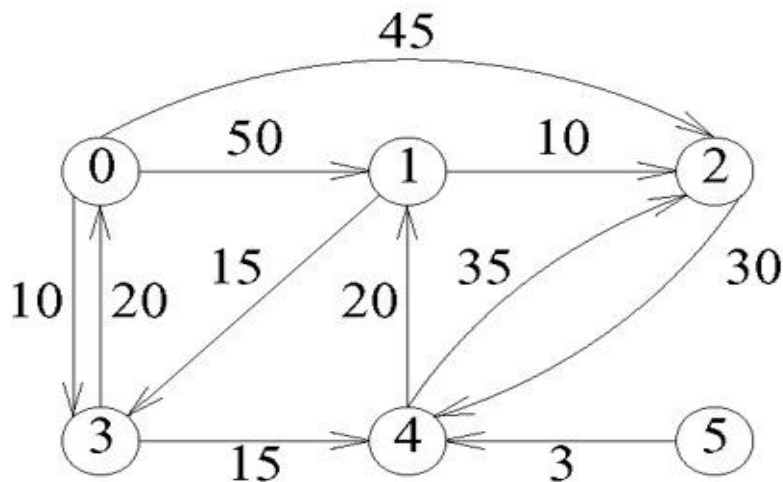
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- Single-Source Shortest Paths
- Algorithm for **Unweighted** Graphs
- Algorithm for **Weighted, Directed Acyclic** Graphs (Weighted DAGs)
- Algorithm for **Weighted, Directed** Graphs with **no negative weights**

# Single-Source Shortest Paths

Give a graph  $G = (V, E)$ , we want to find a shortest path from a given source vertex  $s \in V$  to each vertex  $v \in V$



(a) Graph

<i>Path</i>	<i>Length</i>
1) 0, 3	10
2) 0, 3, 4	25
3) 0, 3, 4, 1	45
4) 0, 2	45

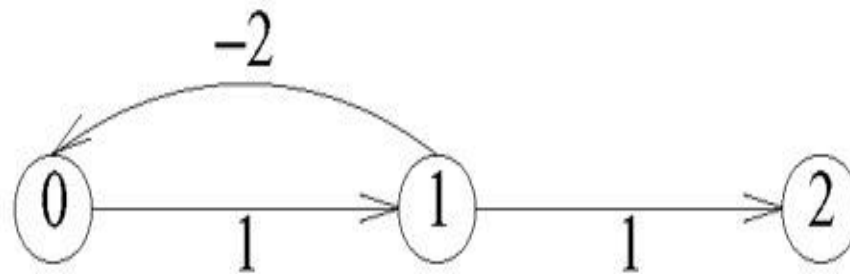
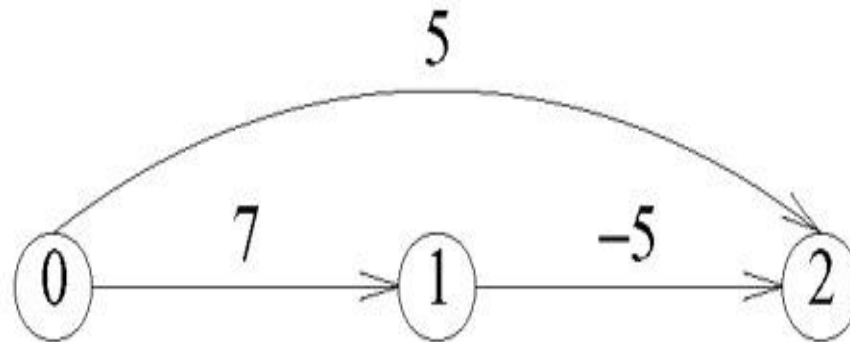
(b) Shortest paths from 0

Why is vertex 5 not included in the list?

# Single-Source Shortest Paths

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What problems might occur with these two special cases?

# Properties of Shortest Paths

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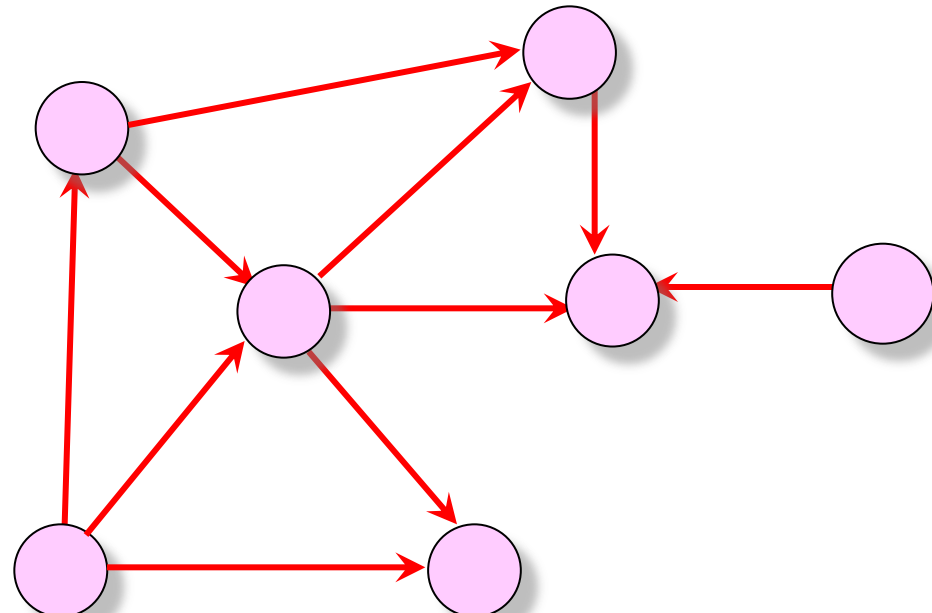
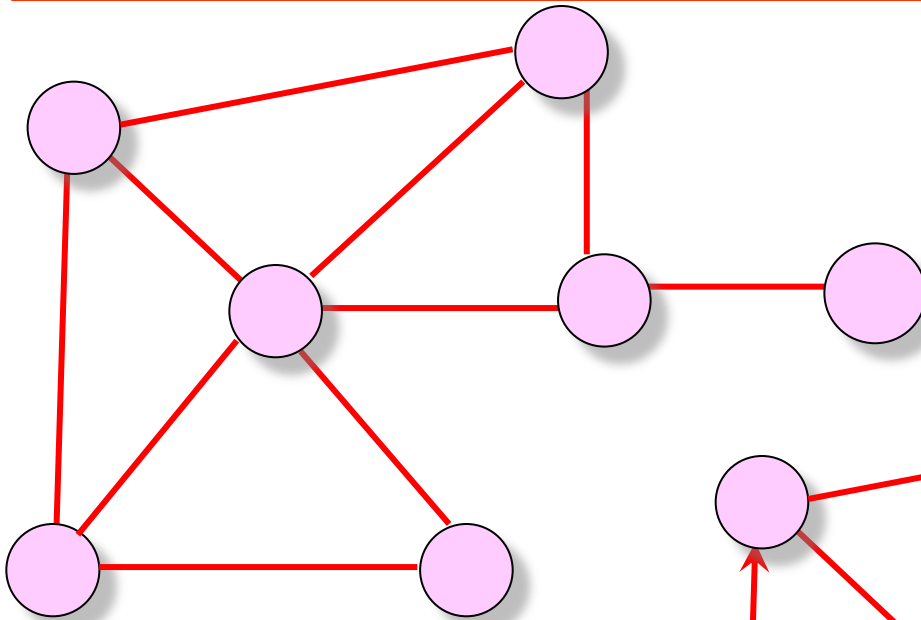
Can a shortest path contain a cycle?

At most how many edges will a shortest path contain?

# Unweighted Graphs

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What algorithm can be used to find the shortest paths for unweighted graphs?

# Shortest Paths Algorithms for Weighted Graphs

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For each vertex  $v \in V$ , we maintain attributes:

$d[v]$  : shortest-path estimate (upper bound on weight of shortest path from source  $s$  to  $v$ )

$\pi[v]$  : predecessor of  $v$  on shortest path so far

Initially  $d[v]$  is  $\infty$  and  $\pi[v]$  is NIL :

INITIALIZE-SINGLE-SOURCE( $G, s$ )

for each vertex  $v \in V[G]$

$d[v] \leftarrow \infty$

$\pi[v] \leftarrow \text{NIL}$

$d[s] \leftarrow 0$

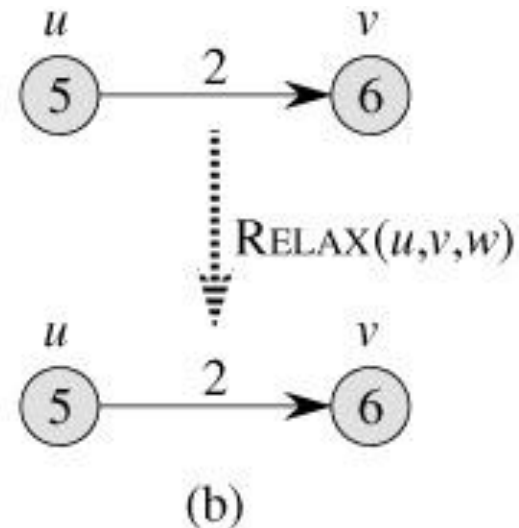
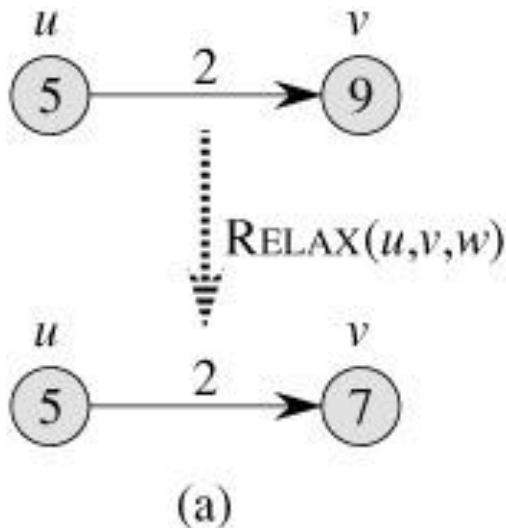
# Updating Adjacent Vertices

RELAX( $u, v, w$ )

if  $d[v] > d[u] + w(u, v)$

$d[v] \leftarrow d[u] + w(u, v)$

$\pi[v] \leftarrow u$





# Algorithm for Weighted DAGs

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AG-SHORTEST-PATHS( $G, w, s$ )

topologically sort the vertices of  $G$

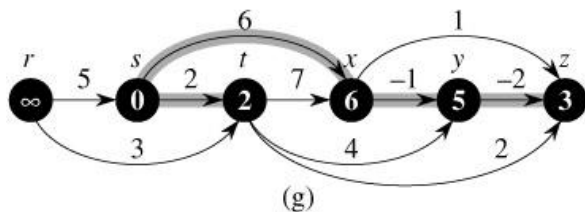
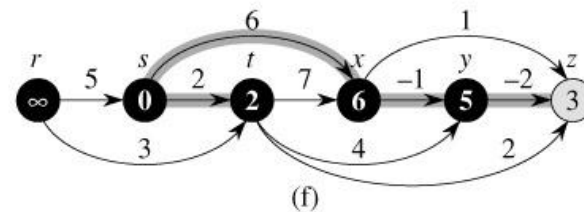
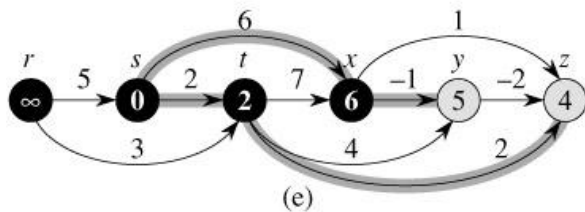
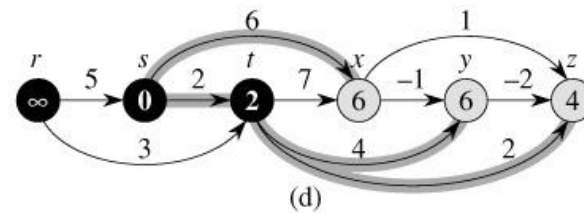
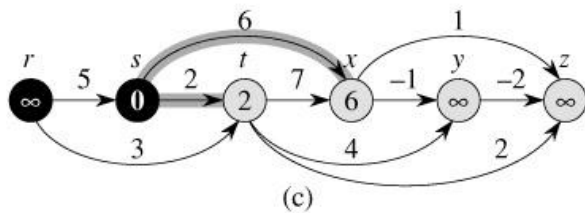
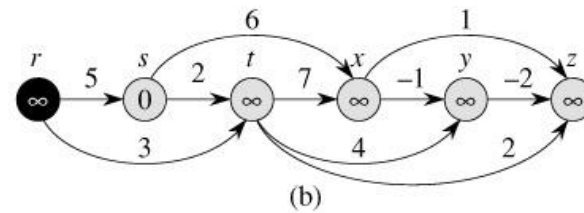
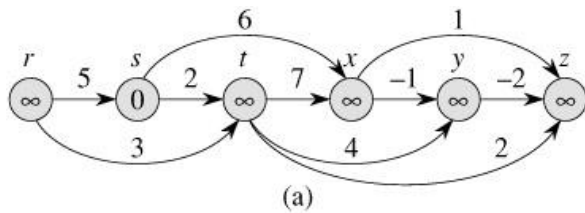
INITIALIZE-SINGLE-SOURCE( $G, s$ )

for each vertex  $u$ , taken in topologically sorted order

for each vertex  $v \in \text{Adj}[u]$

RELAX( $u, v, w$ )

# Example



Shortest paths are always well-defined in dags. Why?

# Weighted Digraphs with No Negative Weights

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# Dijkstra's Algorithm

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DIJKSTRA( $G, w, s$ )

    INITIALIZE-SINGLE-SOURCE( $G, s$ )

$S \leftarrow \emptyset$

$Q \leftarrow V[G]$

**while**  $Q \neq \emptyset$

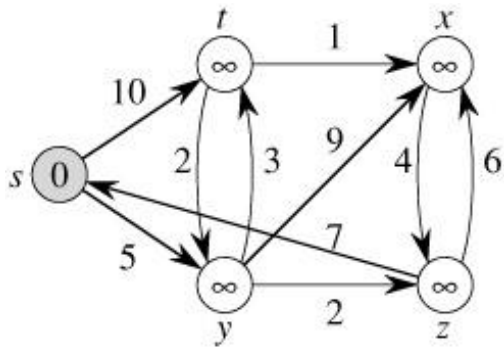
$u \leftarrow \text{EXTRACT-MIN}(Q)$

$S \leftarrow S \cup \{u\}$

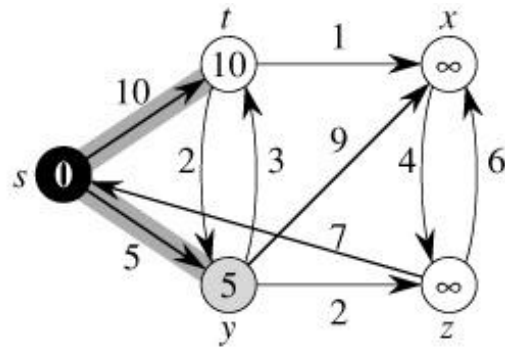
**for** each vertex  $v \in \text{Adj}[u]$

            RELAX( $u, v, w$ )

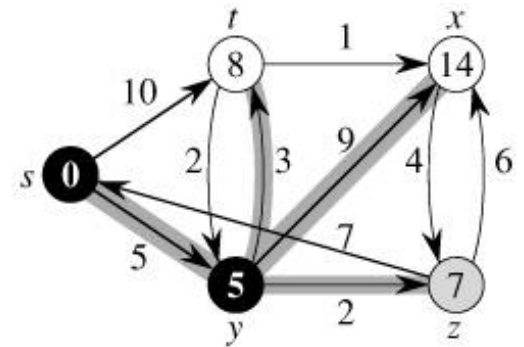
# Example



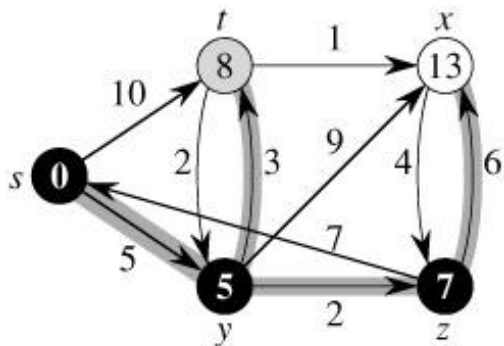
(a)



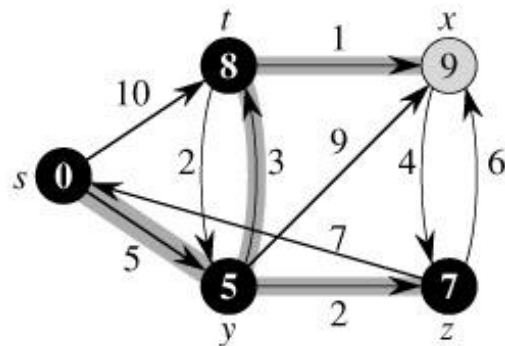
(b)



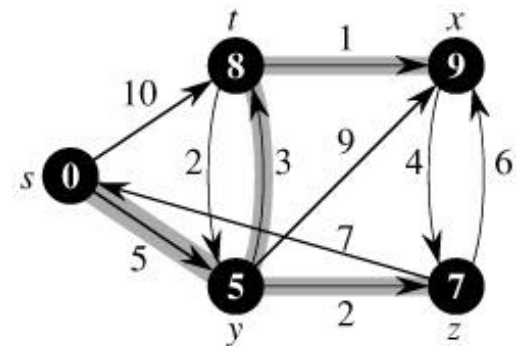
(c)



(d)



(e)



(f)